## Point-to-Subspace Query in $\ell$

- Problem Statement: Given $n$ linear subspaces $\mathcal{S}_{1}, \ldots, \mathcal{S}_{n}$ of $\mathbb{R}^{D}$ of dimension $r$ and a query point $\mathbf{q} \in \mathbb{R}^{D}$, determine the nearest $\mathcal{S}_{i}$ to $\mathbf{q}$ in $\ell^{1}$ norm.


## - Motivation

- Low-dimensional structures in visual data (e.g., lighting, poses) - Structure query as recognition, $\ell^{1}$ for robustness (to, e.g., occlusions, shadows)
Efficiency: large $D$ (e.g., \# pixels) and large $n$ (many subjects)


## Existing Work

Efficiency: Preprocessing and storage - lowpoly ( $D, r, n$ ), Query - lowpoly ( $D, r, n^{o(1)}$ )
$r=0$
$\ell^{2} \checkmark($ e.g., LSH $)$
$\checkmark$ (Andoni et al, SODA'09) $\quad \begin{array}{r}r \geq 2 \\ \text { ? Heuristics (Bastri'11, Jair'10 }\end{array}$


- © Sublinear time algorithm for point-to-hyperplane in $\ell^{2}$ (and $\ell^{1}$ ) unlikely (Williams'05)
$\bullet$ General low-distortion low-dimensional embedding for $\ell^{1}$ impossible (Brinkman'05)
- Precursor $:$ : For a single subspace, $\ell_{\square}^{1} \rightarrow \ell_{1}^{(\text {(rlogr) }}$ with distortion $O(r \log r)$ (Sohler and Woodruff, 11)
- Error-Correction in $\ell^{1}$ : dim. reduction determined by density of error $e$ (Candes and Tao, 04)

Our Algorithm and Main Results


- Cauchy distribution: $p(x)=\frac{1}{\pi 1+x^{2}}$
- No finite mean or variance
- $\ell^{1}$ Stable: for iid standard Cauchy RV's $\phi$ $\sum_{i=1}^{k} \phi_{k} \sim\|\Phi\|_{\ell^{1}} x$, for $x$ standard Cauchy.
- Algorithm: Generate a random matrix $\mathbf{P} \in \mathbb{R}^{d \times D}$ with iid Cauchy RV's $(d \ll D)$

$$
\text { Preprocessing: Compute the projections } \mathbf{P} \mathcal{S}_{1}, \cdots, \mathcal{P} \mathcal{S}_{r}
$$

$$
\text { Test: Compute the projection } \mathbf{P q} \text {, and compute its } \ell^{1} \text { distance to each of } \mathbf{P} \mathcal{S}_{i}
$$

- Theory: In short, Cauchy projection with large enough $d$ preserves the identity of nearest subspace with nontrivial probability. In full details

Suppose we are given $n$ linear subspaces $\left\{\mathcal{S}_{1}, \cdots, \mathcal{S}_{n}\right\}$ of dimension $r$ in $\mathbb{R}^{D}$ and any query point $\mathbf{q}$, and that the $\ell^{1}$ distances of $\mathbf{q}$ to each of $\left\{\mathcal{S}_{1}, \cdots, \mathcal{S}_{n}\right\}$ are $\xi_{1} \leq \cdots \leq \xi_{n^{\prime}}$ when arranged in ascending order, with $\xi_{22} / \xi_{1} \geq \eta>1$. For any fixed $\alpha<1-1 / \eta$, there exists $d \sim O(r \log n)^{1 /}$ (assuming $n>r$ ), if $\mathbf{P} \in \mathbb{R}^{d \times D}$ is iid Cauchy, we have

$$
\underset{i \in[n]}{\arg \min } d_{\ell^{\prime}}\left(\mathbf{P q}, \mathbf{P} \mathcal{S}_{i}\right)=\underset{i \in[n]}{\arg \min } d_{\ell^{\prime}}\left(\mathbf{q}, \mathcal{S}_{i}\right)
$$

with (nonzero) constant probability

- Implications:
- $d$ depends on the relative gap $n$, and not on $D$.
- d depends on $\log n-$ growing nicely wrt. \# subspaces.
- Independent trials can be taken to amplify the success probability. (Matter of low-dimensional $\ell^{1}$ regressions!
- In case of ties, first $k$ nearest neighbors can instead be considered.


## Empirical Results

- Extended Yale B Face Data: $D \sim 30,000, n=38, d=9$ Subset under moderate lighting (single projection)

- Normalized distance gap for moderately/extremely illuminated samples.



## Bounded expansion for the good subspace

- Distance after projection can easily be upper bounded:

$$
d_{\ell^{\prime}}\left(\mathbf{P q}, \mathbf{P} \mathcal{S}_{\star}\right)=\min _{\mathbf{h} \in \mathcal{P}_{\star}}\|\mathbf{P q}-\mathbf{h}\|_{1}
$$

$$
\leq \stackrel{\mathbf{h} \in \mathbf{P} \mathcal{S}_{\star}}{\left\|\mathbf{P q}-\mathbf{P} \mathbf{v}_{\star}\right\|_{1}=\left\|\mathbf{P}\left(\mathbf{q}-\mathbf{v}_{\star}\right)\right\|_{1} . . . . . . .}
$$

- Bounded expansion with nonzero constant probability: There exists numerical constant $c \in(0,1)$ with the following property. If $\mathbf{w} \in \mathbb{R}^{D}$ be any fixed vector, and suppose that $\mathbf{P} \in \mathbb{R}^{d \times D}$ is a matrix with iid standard Cauchy entries. Then for any $\rho>1$,

$$
\mathbb{P}\left[\|\mathbf{P w}\|_{1}>\rho \frac{2}{\pi} d \log d\|\mathbf{w}\|_{1}\right]<c+\frac{1-c}{\rho}<1 .
$$

## Idea of Proof

- Behavior of Cauchy projection on point-to-subspace $\ell^{1}$ distance for different subspace configurations

- P does not increase the distance to "good" subspace too much, and - P does not shrink the distances to "bad" subspace too much.


## Bounded contraction for the bad subspaces

- Consider $\forall \mathbf{w} \in \mathcal{S}_{i} \oplus \mathbf{q}$, want to show $\|\mathbf{P w}\|_{1} \geq \gamma\|\mathbf{w}\|_{1}$ for some appropriate $\gamma$, then $d_{\ell^{1}}\left(\mathbf{P q}, \mathbf{P} \mathcal{S}_{i^{\prime}}\right)=\min _{\mathbf{v} \in \mathcal{S}_{i}}\|\mathbf{P q}-\mathbf{P v}\|_{1} \geq \min _{\mathbf{v} \in \mathcal{S}_{i}}\|\mathbf{P}(\mathbf{q}-\mathbf{v})\|_{1}$ $\geq \min _{\mathbf{v} \in \mathcal{S}_{i}} \gamma\|\mathbf{q}-\mathbf{v}\|_{1}=\gamma d_{\ell^{1}}\left(\mathbf{q}, \mathcal{S}_{i}\right)$,
- Discretization argument on restricted unit $\ell^{1}$ sphere onto the augmented subspaces $\Gamma=\left\{\mathbf{w} \mid\|\mathbf{w}\|_{1}=1\right\} \cap \tilde{\mathcal{S}}_{j}$. Points on the $\varepsilon$-net covered by the concentration results on lower tail. Let $\mathbf{P} \in \mathbb{R}^{d \times D}$ be an iid Cauchy matrix. Then for any fixed vector $\mathbf{w} \in \mathbb{R}^{D}$ and $\alpha, \delta \in(0,1)$,

$$
\mathbb{P}\left[\|\mathbf{P w}\|_{1}<(1-\alpha)(1-\delta) \frac{2}{\pi} d \log d\|\mathbf{w}\|_{1}\right]<d^{1-\alpha} \exp \left(-\frac{\delta^{2}}{2 \pi} d^{\alpha}\right)
$$

- Points on 「 but off the $\varepsilon$-net covered by triangular inequalty, which is founded on well-conditioned basis for $\ell^{1}$ subspaces.

